**CSE 5350/7350 Due Date:**

**D.W. Matula 04 Oct. 2011**

**Homework Set # 2**

**Data Structure Engineering for Algorithm Design**

**Data Structures** deals with the **representation** and **organization** of data to support efficient:

        **storage**

        **access**

        **execution of operations**

        **display** **of results**.

1.  **[Data structure conversion](15 points)**. Design and analyze an on-the-fly algorithm for converting a binary search tree to an ordered doubly linked list following the inorder traversal of the data. Utilize a loop invariant (text p.17) to argue the correctness of your algorithm. Have your algorithm use as little extra space as possible, with "in place" the goal. Perform a walkthrough for the sequence S where the elements of S are first processed left-to-right to build the search tree. Then apply your algorithm to convert the binary search tree to a doubly linked list, showing the status of the conversion at each node of the search as it is consumed and moved into the doubly linked list.

S = 54, 78, 63, 35, 77, 49, 83, 28, 14, 45, 70, 12.

2.  **[Heapify] (15 points).** The array form of a heap (i.e. a binary heap array) has the children of the element at position *i* in positions (left @ 2*i*, right @ 2*i*+1). For a balanced ternary heap array the three children of the element at position *i* are at positions (left @ 3*i*-1, middle @ 3*i*, right @ 3*i*+1). The term "Heapify" refers to the algorithm that builds a binary heap array from right-to-left (leaf-to-root in binary tree form) given the size of the array to be built. The term "Ternary Heapify" here refers to the analogous algorithm that builds a balanced ternary heap array in right-to-left order.

a. Heapify the sequence S of Problem 1, and give the number of comparisons required.

b. Do the same as (a) for Ternary Heapify.

c. Build the decision tree determined by the comparisons employed to Heapify a six element list: a1, a2, a3, a4, a5, a6.What is then the worst case and average case number of comparisons to Heapify six elements? Is this optimal in either case?

d. Do the same as (c) for Ternary Heapify.

3.  **[Graph (Symmetric Matrix) Reordering] (15 points).** Let the adjacency list representation of a graph ( 0, 1 symmetric matrix ) be given by a packed edge list. Describe an algorithm for reordering the edge list in place so all adjacency lists follow a new vertex ordering given by a permutation p(1), p(2), ..., p(n). Apply your algorithm to the graph given by the packed edge list stored as example HW2 - 3 graph. Reorder by determining a maximum adjacency search order(see Project 1(a) on class homepage) with "ties" broken lexicographically.

HW2 - 3 graph in adjacency list form:

* + a: c, f, g, h;
  + b: c, d;
  + c: a, b, g, h;
  + d: b, e, f, h;
  + e: d, f, h;
  + f: a, d, e, g;
  + g: a, c, f;
  + h: a, c, d, e

4. **[Radix and Digit Set Conversion]**

Note: Explain your conversion procedures for each part. (ref: pp 59-65 of K+M).

1. Convert the following 16-bit integer *x* to standard radix 4 and to Booth recoded radix 4 (digit set {-2, -1, 0, 1, 2}). Check the answers by determining digit sums module 3.

*x* = 1010 0101 1100 01112

1. Convert the following 6-digit decimal *y* to balanced ternary digit serially low order digit first using decimal (source system) arithmetic, and then high order digit first using balanced ternary (target system) arithmetic.

*y* = 314, 159

5. **[Sparse matrix multiplication](15 points).** Describe an efficient sparse matrix multiplication procedure for two n x n matrices stored in "adjacency list" form. Assume matrix A has mA non zero entries and matrix B has mB non zero entries, with n << mA << mB << n2. Describe why your algorithm will generate the n × n matrix product in adjacency list form in time O (n mA).

Apply your algorithm for multiplying the two 8 × 8 0,1 matrices A1, B1 where A1 is given by the following adjacency list of unit entries and B1 is given by the adjacency list of unit entries of Problem 3. Count the actual number of multiplications used in this matrix product and explain why it is considerably less than the bound n mA (which is 8 × 19 =152) in this case.

*Adjacency List For 0,1 Matrix A1*

* + 1: 3, 5, 7
  + 2: 4, 8
  + 3: 5, 8
  + 4: 2, 6
  + 5: 1, 2, 7, 8
  + 6: 3
  + 7: 1, 5, 6
  + 8: 6, 7