

CSE 2353 SPRING 2008
DISCRETE COMPUTATIONAL
STRUCTURES

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CSE2353 SPRING 2008 OUTLINE

- PART ONE
 - I. Introduction
 - II. Sets
 - III. Logic

- PART TWO
 - IV Proof Techniques
 - V. Counting Principles & Combinatorics
 - VI. Relations & Functions

- PART THREE
 - VII. Graphs & Trees
 - VIII. Boolean Algebra

Note: All book references are to *Discrete Mathematics with Proof* by Eric Gossett

CSE2353 - INTRODUCTION

- What is Discrete?
- “Consisting of distinct or unconnected elements.”
Webster’s Seventh New Collegiate Dictionary
- Course objectives: Study discrete math concepts which are the basis of programming and many Computer Science topics
- Will relate each topic to Computer Science

Why discrete?

CSE2353 - SETS OUTLINE

- Basic Definitions
- Venn Diagrams
- Set Operations
- Power Set
- Relationship to Computer Science

SET DEFINITIONS

- Well defined collection of distinct objects.
- Set $\{ \dots \}$
- Element of \in
- Cardinality $| \dots |$
- Subset \subset, \subseteq
- Equality p16

DEFINING A SET

- Enumerate (Roster Method)
- Example: $A = \{1, 2, 3\}$
- Identify set membership (Set Builder)
- Example: $A = \{x \mid x \in \mathbb{Z} \text{ and } x \geq 1 \text{ and } x \leq 3\}$

SPECIAL SETS

- Null (Empty) Set - \emptyset
- Universal Set - U
- Natural Numbers - \mathbf{N}
- Integers - \mathbf{Z}
- Positive Integers - \mathbf{Z}^+
- Negative Integers - \mathbf{Z}^-
- Rational - \mathbf{Q}
- Real - \mathbf{R}
- Complex - \mathbf{C}
- Power - $\mathbf{P}(A)$

If $|A| = n$, what is $|\mathbf{P}(A)|$?

SET OPERATIONS

- Intersection \cap
- Union \cup
- Difference -
- Complement \bar{A}
- We can represent these using Venn Diagrams
- Cartesian Product \times
- Laws on p 27

Venn Diagram Worksheet

SETS IN COMPUTER SCIENCE

- Representation:
 - Linked List
 - Array
 - Bit Map
- Multiset (Removing duplicates is sometimes difficult)
- Ordered Set
- File
 - Subset of Cartesian Product
 - Operations can be viewed as set operations

CSE2353 - LOGIC OUTLINE

- Propositions
- Truth Tables
- Predicates
- Use in CS

WHAT IS A PROPOSITION?

- Proposition - Statement which is true or false.
- Examples
 - There are planets circling Alpha Centauri.
 - 3 is an even
- Propositional Calculus - Notation and rules for combining and manipulating propositions.
- Notation:
 - Proposition: p, q, r, s, \dots (Small or Capital)
 - Operations:
 - * and (conjunction) \wedge
 - * or (inclusive disjunction) \vee
 - * not (negation) \neg or \sim or \bar{p}
 - * if ... then (implication) \rightarrow or \supset
 - * if and only if (biconditional) \leftrightarrow
 - Precedence: $\neg \wedge \vee \rightarrow \leftrightarrow$

TRUTH TABLES

- Truth Tables - Summarizes truth values of propositions.
- Column for each proposition (simple and compound)
- Row for each possible combination of truth values for simple propositions
- See tables on p34, p35

Examples: \bar{p} , $p \wedge q$, $p \vee q$, ...

IMPLICATION

- Examples
 - If 3 is even then 4 is odd.
 - If I get an A in this course I'll celebrate by going to Europe
- $p \rightarrow q$
- p - antecedent (hypothesis)
- q - consequent (conclusion)
- Converse of $p \rightarrow q$ is $q \rightarrow p$
- Inverse of $p \rightarrow q$ is $\bar{p} \rightarrow \bar{q}$
- Contrapositive of $p \rightarrow q$ is $\bar{q} \rightarrow \bar{p}$

Truth Table for Implication

LOGICAL EXPRESSIONS

- Well Formed Formula:
 - $\langle prop \rangle ::= p \mid q \mid \dots \mid z$
 - $\langle binop \rangle ::= \rightarrow \mid \wedge \mid \vee \mid \leftrightarrow$
 - $\langle wff \rangle ::= \langle prop \rangle \mid (\overline{\langle wff \rangle}) \mid (\langle wff \rangle \langle binop \rangle \langle wff \rangle)$
- Note: Parentheses may be dropped if order of execution follows precedence order

How many different logical expressions are there on two propositions?

OTHER OPERATIONS

- xor \oplus ($\underline{\vee}$)
- nand $|$
- nor \downarrow
- tautology (denoted by t or T)
- contradiction (denoted by f or F)

LOGIC WORKSHEET

EQUIVALENCE and INFERENCE

- Two propositions are logically equivalent ($p \equiv q$) or ($p \Leftrightarrow q$) if they have the same truth values ($p \leftrightarrow q$ is T).
- A proposition logically implies (infers) another proposition ($p \vdash q$) or ($p \Rightarrow q$) if whenever p is true then q is also true (Note that $p \rightarrow q$ is T.)

How to show equivalence?

How to show inferences?

INFERENCE RULES FOR PROPOSITIONAL CALCULUS

- Inference 1: (Modus Ponens)
 - From $p \rightarrow q$ and p , infer q .

- Inference 2: (Substitution)
 - A propositional variable may be substituted by another variable.

- With substitution must be careful to substitute all occurrences of a propositional variable.

- Substitution p45

- Tables 2.21, 2.22, 2.23 p47-48

Equivalence/Logical Inference Worksheet

ARGUMENT

- Argument (Claim)
 - Set of propositions (premises) and another proposition (conclusion).
 - The argument is valid if the conjunction of the premises logically implies the conclusion. Else it is invalid.

- Notation
 - $A_1, A_2, \dots, A_{n-1} \Rightarrow A_n$
 - A_1, A_2, \dots, A_{n-1} are the premises
 - A_n is the conclusion
 - This is a valid argument if $A_1 \wedge A_2 \wedge \dots \wedge A_{n-1} \rightarrow A_n$ is T

- Section 2.7.1 p69-72

SYLLOGISMS

- Techniques to infer new true statements from existing statements
- Major premise, Minor premise, Conclusions
- Syllogisms p14
- Valid arguments depend on the application of syllogisms
- <http://en.wikipedia.org/wiki/Syllogisms>
- Disjunctive Syllogism
- Hypothetical Syllogism

Valid Argument Worksheet

WHAT IS A PREDICATE?

- Not all statements are propositions
 - $2x$ is even
 - $x+y=4$
- However these can be expanded to become propositions
- Predicate - property describing an object.
- Predicates are denoted by capital letters: E, F
- Predicate Form (Propositional Function) - Predicate with variables: $E(x), F(x,y)$
- A predicate form becomes a proposition when variables are replaced with elements from a specific domain: $E(2), F(2,3)$

Can you determine the truth value of a predicate. Why?

What do you have to do to convert a predicate to a proposition?

QUANTIFIERS

- Can also convert a propositional function into a proposition by using quantifiers, \forall, \exists :
 - $\forall x, E(x)$
 - $\exists x, \exists y, F(x,y)$
- Examples
 - Computer Science students are brilliant: $B(x)$
 - All computer science students are brilliant:
 $\forall x, B(x)$
 - Some computer science students are brilliant:
 $\exists x, B(x)$

Quantifier Worksheet

ARGUMENTS IN PREDICATE LOGIC

- Claims with quantifiers, p73
- Universal Specification: If $\forall x, F(x)$ is true then $F(a)$ is true for all a in the universe of discourse
- Universal Generalization: If $F(a)$ is true for all a in the universe of discourse then $\forall x, F(x)$ is true.
- Existential Specification: If $\exists x, F(x)$ is true then there is at least one a in the universe of discourse such that $F(a)$ is true
- Existential Generalization: If $F(a)$ is true for some a in the universe of discourse then $\exists x, F(x)$ is true.

Inference Rules Handout**Examples**

RELATIONSHIP OF LOGIC TO PROGRAMMING

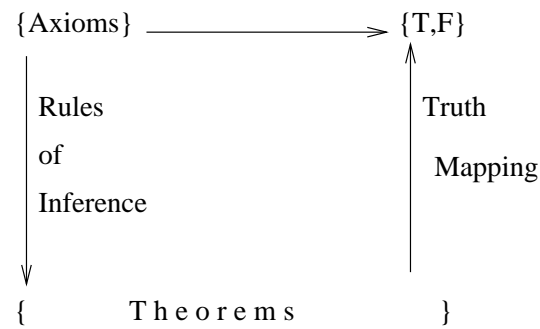
- Conditional statements
- Variables
- Looping
- Database query languages

CSE2353 - PROOF TECHNIQUES OUTLINE

- What is a proof?
- Proof Techniques
- Examples - practice
- Relationship to Computer Science

AXIOM SYSTEM

- Axiom - Statement which is assumed. Given as a starting point.
- Inference Rules - Rules describing how to generate new statements from existing ones.
- Theorem - Statement derivable from a set of axioms using specified rules of inference.
- Proof - Finite sequence of statements where each is either an axiom or theorem previously derived.



PROVING A THEOREM

- A theorem is a statement which is logically implied by the conjunction of axioms.
- A theorem is a proposition T , such that
 - $(A1 \wedge A2 \wedge \dots \wedge An) \vdash T$
- Assume the truth of the axioms and show this guarantees the truth of T
- $P \Rightarrow Q$ or P implies Q . This means that Q follows from P . Or, the conjunction of P and all axioms logically implies Q .

METHODS OF PROOF

- Vacuous
- Trivial
- Direct
- Contrapositive
- Contradiction
- Induction
- YOU CAN NOT PROVE BY EXAMPLE
- Counter Example
- Divide into Cases
- Construction

IMPLICATION "PROOFS"

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

- Vacuous - Only rows 1 and 2 occur. Show p is false
- Trivial - Only rows 2 and 4 occur. Show q is true
- Direct - Only rows 1,2,4 occur
- Indirect/Contrapositive - Prove $\neg q \rightarrow \neg p$.
- Contradiction - Show $\neg(p \wedge \neg q)$ is a tautology.
- These show the existence of a proof but are not technically a proof.

Proof Handout

EXAMPLE PROOFS

Prove: If $x \in \mathbb{Z}$ is even then $x+1$ is odd.

Prove: There is no largest prime number.

Prove: If A and B are finite sets, then $A \cup B$ and $A \cap B$ are also finite.

Direct Proofs Worksheet

Indirect Proofs Worksheet

Contradiction Proofs Worksheet

PROOF BY INDUCTION

- Induction allows us to prove a property about an infinite (or a large) set of objects without having to look at all the objects.
- Ex: I can climb to the top of the Empire State Building.
- Technique:
 - Prove conjecture holds for $n=1$. (Basis/Base)
 - Prove that for all $k \geq 1$, if the result holds for $n=k$, then it must also hold for $n=k+1$ (Inductive step: Hypothesis and Induction).

Prove

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

Proof by Induction Worksheet

General Proofs Worksheet

CSE2353 COUNTING PRINCIPLES & COMBINATORICS OUTLINE

- Basic Counting Principles Addition and Multiplication
- Pigeonhole Principle
- Permutations and Combinations

Addition & Multiplication Principles

- Suppose tasks T_1, T_2, \dots, T_n can be done in x_1, x_2, \dots, x_n ways.
- Addition Principle Number of ways to do one
 - The number of ways any one of the tasks can be done is $x_1 + x_2 + \dots + x_n$.
 - Example: Suppose a menu has three different desserts: cake, ice cream and cookies. If there are 5 types of ice cream, 3 types of cakes, and 10 types of cookies, how many actual choices does one have to choose a dessert?
- Multiplication Principle Number of ways to do all
 - The number of ways any all of the tasks can be done is $x_1 \cdot x_2 \cdot \dots \cdot x_n$.
 - Example: How many ways could one order one of each type of dessert?

Counting Principles Worksheet

Pigeonhole Principle

- Suppose there are n pigeons and k pigeonholes where $n > k$. If all these pigeons try to fit into the pigeonholes, then some pigeonholes will have to contain at least two pigeons.
- Generalized Pigeonhole Principle: n pigeons, k pigeonholes, where $m = \lceil \frac{n}{k} \rceil$. Some pigeonhole must contain at least m pigeons.
- Example: Suppose a room has 25 people in it? Then at least 3 must have a birthday in the same month.

Pigeonhole Principle Worksheet

Permutations and Combinations

- Permutations: Number of ways to arrange r objects from a set of n objects in order (without repetition) is: $P(n, r) = \frac{n!}{(n-r)!}$
- Combinations Number of ways to select r objects from a set of n objects (order independent, no repetition): $C(n, r) = \frac{n!}{(n-r)! \cdot r!}$
- How many different combinations of desserts could a table of 4 have? (Use the dessert choices mentioned earlier.)

Permutations and Combinations Worksheet

CSE2353 - RELATIONS & FUNCTIONS OUTLINE

- What is a relation?
- Equivalence relation
- Partial order
- Relations in Computer Science
- What is a function?
- Function operations
- Functions in Computer Science

RELATION DEFINITIONS

- Relation
 - $R \subseteq A \times B$.
 - $x \mathbf{R} y$ iff $(x, y) \in \mathbf{R}$.
 - A is the domain of the relation
 - B is the range of the relation
 - NOTE: Some definitions of range/domain only refer to the subset of B/A actually involved in the relation
- Tuple - element of a relation.
- Relation on a set A , is a subset of $A \times A$.
- Inverse Relations - $x \mathbf{R}^{-1} y$ iff $y \mathbf{R} x$.

REPRESENTING BINARY RELATIONS

- Set - ordered pairs
- Set definition - membership values
- Arrow Diagram
 - Write elements of set A in one column and elements of B in another
 - Draw an arrow from each element in A column to each element in B column with which it is related
- Digraph (Directed Graph) for relation on set A
 - Create a vertex for each element in A and label
 - Draw an arrow from each vertex in A to each vertex with which it is related

Relation Worksheet

PROPERTIES OF RELATIONS

- Let \mathbf{R} be a relation on a set A :
 - Reflexive: $\forall a \in A \ a \mathbf{R} a$.
 - Symmetric: $\forall a \in A$ if $a \mathbf{R} b$ then $b \mathbf{R} a$.
 - Antisymmetric: $\forall a \in A$ if $a \mathbf{R} b$ and $b \mathbf{R} a$ then $a=b$.
 - Transitive: $\forall a \in A$ if $a \mathbf{R} b$ and $b \mathbf{R} c$ then $a \mathbf{R} c$.
 - Irreflexive: $\neg \exists a \in A$ where $a \mathbf{R} a$.

- Examples:
 - Reflexive: $\subseteq, \leq, =$.
 - Symmetric: $=$.
 - Antisymmetric: $\subseteq, \leq, =, <$
 - Transitive: $\subseteq, \leq, =$.
 - Irreflexive: $<, >$.

- The closure of a relation \mathbf{R} , is the smallest superset of R which satisfies a particular property.

Relation Properties Worksheet

EQUIVALENCE RELATIONS

- Equivalence - reflexive, symmetric, transitive
 - Examples: $=$, mod n
 - Partition of a set S is a collection of subsets $A_1, A_2, \dots, A_n, \dots$ such that $A_1 \cup A_2 \cup \dots \cup A_n \cup \dots = S$ and if $i \neq j$ then $A_i \cap A_j = \emptyset$.
 - An equivalence relation on a set S defines a set of equivalence classes which partition S .
 - The equivalence class of $x \in A$ is the set of elements of A to which x is related: $[x]_R = \{y \in A \mid x R y\}$.

PARTIAL ORDER RELATIONS

- Partial Order - reflexive, antisymmetric, transitive
 - Examples: \leq , \subseteq
 - Linear Order (Total Order) - partial order where any two sets are related
 - Poset - set with a partial order.
 - Hasse Diagram - technique of representing a finite poset.

MORE ON POSETS

- Given A, \mathbf{R} which is a partial order:
 - Greatest element of A : $\alpha \in A$ such that $\forall a \in A \ a \mathbf{R} \alpha$
 - Least element of A : $\beta \in A$ such that $\forall a \in A \ \beta \mathbf{R} a$
 - Poset need not have a greatest or least element
 - Maximal element of A : $x \in A$ if $\forall a \in A \ x \mathbf{R} a$ implies $x=a$
 - Minimal element of A : $y \in A$ if $\forall a \in A \ a \mathbf{R} y$ implies $a=y$

Poset and Equivalence Relations Worksheet

RELATIONS IN COMPUTER SCIENCE

- Relational Databases (Sometimes called a table)
- Relational Algebra
- Relational Calculus

Relational database worksheet

FUNCTION DEFINITIONS

- Function - Function f from A to B , $f:A\rightarrow B$, is a relation that associates to each $a\in A$ a unique element $f(a)\in B$. (Also called mapping or transformation)
- A is the domain
- B is the codomain
- Vertical line test
- Range of f (or image) - $\text{im}(f) = \{f(a):a\in A\}$
- Identity function - $\text{id}_A:A\rightarrow A$ where $\text{id}_A = \{(x,x):x\in A\}$

COMPOSITE FUNCTIONS

- Let $f:A\rightarrow B$ and $g:B\rightarrow C$ be functions. The composite function $g\circ f:A\rightarrow C$ is

$$g\circ f = \{(x,z) \in A \times C : (x,y) \in f \text{ and } (y,z) \in g \text{ for some } y \in B\}$$

- $g\circ f(x) = g(f(x))$
- $\text{im}(g\circ f) \subseteq \text{im}(g)$

PROPERTIES OF FUNCTIONS

- Let $f:A \rightarrow B$ be a function:
 - Injection(1-1) - $\forall a, a' \in A, \text{ if } a \neq a'$
then $f(a) \neq f(a')$
 - Surjection(Onto) - $\forall b \in B, \exists a \in A$
such that $(a, b) \in f$
 - Bijection - both 1-1 and onto
 - If f is a bijection then $g:B \rightarrow A$ defined by
 $g(b)=a$ iff $f(a)=b$ is the inverse of f and is
denoted by f^{-1}
- The composite of injective (surjective) functions
is injective (surjective)
- If a composite $g \circ f$ is injective (surjective) then so
is f

Functions Worksheet

CARDINALITY

- Two sets have the same cardinality if there is a bijection between them.
- Countably Infinite sets have a bijection from the positive integers

APPLICATIONS OF FUNCTIONS IN COMPUTER SCIENCE

- Functions/Procedures in Programming Languages
- Functional Dependency
- Complexity Analysis

CSE2353 - GRAPHS & TREES OUTLINE

- What is a Graph?
- Special Types of Graphs
- Trees
- Fun Graph Problems
- Graphs in Computer Science

WHAT IS A GRAPH?

- An (Undirected) Graph G consists of:
 - A finite non-empty set V of Vertices
 - A finite set E of Edges
 - A function $\delta: E \rightarrow P(V)$ such that, for every edge e , $\sigma(e)$ is a one- or two-element subset of V .

GRAPH TERMS

- Adjacent - Vertices which are adjacent if they have an edge in common. Edges are adjacent if they have at least one vertex in common.
- Incident - Two vertices are incident to an edge if they are connected via that edge. In this case the edge is incident to the two vertices.
- Degree - The degree of a vertex is the number of edges which are incident to it.
- Subgraph - H is a subgraph of G if its vertices and edges are subsets of those in G .
- Edge Sequence - Sequence of edges such that any two consecutive ones are adjacent.
- Path - Edge sequence in which all edges are distinct
- Circuit - Closed (starting and ending vertex the same) simple (all vertices distinct) path with at least one edge.

GRAPH REPRESENTATIONS

- Suppose $G = \{V, E\}$ where $|V| = n$ and $|E| = m$
- Adjacency Matrix $n \times n$ matrix A where $A_{i,j}$ is the number of edges joining v_i and v_j . (Note if only one edge exists, the value may be a weight associated to that edge.)
- Incidence Matrix $m \times n$ binary matrix B where $B_{i,j}$ is 1 iff e_i is incident with v_j .

DIRECTED GRAPHS

- Graph where edges have a direction
- Edges are sometimes called arcs

Graphs Worksheet

SPECIAL GRAPHS

- Simple - No loops or multiple edges
- Null - No edges.
- Complete K_n - Simple graph where every pair of distinct vertices is joined by an edge.
- Bipartite - Vertices can be partitioned such that every edge joins a vertex from one partition to that in another partition.
- Connected - Any pair of distinct vertices are connected via a path.

TREES

- Tree - Connected graph with no circuits
- Rooted Tree
 - Root
 - Leaves
 - Parent
 - Child
- Binary Tree - Every vertex has cardinality 2 or 3 (0 or 1 children)

Trees Worksheet

FUN GRAPH PROBLEMS

- Königsberg Bridge Problem p8
- Traveling Salesman Problem
- Stable Marriage Problem p3
- Four Color Problem

GRAPHS IN COMPUTER SCIENCE

- Graph Representation
 - Array
 - Linked List
 - Tree - One dimensional array
- Algorithms
- Applications
 - Indexing Data Structures (B⁺ Tree)
 - Spelling Checkers (Trie)
 - Game Trees
 - Efficient Routing
 - Parsing
 - Sorting
 - Scheduling programs - Heap
 - Hierarchical and Network Databases
- NOTE - Graph proofs often use Induction

Finite State Machine

CSE2352 - BOOLEAN ALGEBRA OUTLINE

- Algebraic Structures
- What is a Boolean Algebra?
- Relationship to Sets and Propositions
- Boolean Functions
- Minterms/Maxterms
- Boolean Algebra in Computer Science

WHAT IS AN ALGEBRAIC STRUCTURE?

- Algebraic Structure - One or more sets with one or more operations
- Example: \mathbf{Z} with $+, -, *$
- Example: Given S , \cap , \cup , $-$ on $P(S)$
- Algebraic Structures are classified based on the properties of the operations
- The primary Algebraic Structure we will look at is a Boolean Algebra

BOOLEAN ALGEBRA INTRODUCTION

 $(\{0,1\}, +, *, ^-, 0, 1)$

$+$	0	1
0	0	1
1	1	1

$*$	0	1
0	0	0
1	0	1

b	\bar{b}
0	1
1	0

 $(P(\{x\}), \cup, \cap, ^-, \emptyset, \{x\})$

\cup	\emptyset	$\{x\}$
\emptyset	\emptyset	$\{x\}$
$\{x\}$	$\{x\}$	$\{x\}$

\cap	\emptyset	$\{x\}$
\emptyset	\emptyset	\emptyset
$\{x\}$	\emptyset	$\{x\}$

b	\bar{b}
\emptyset	$\{x\}$
$\{x\}$	\emptyset

 $(\{f, t\}, \vee, \wedge, ^-, f, t)$

\vee	f	t
f	f	t
t	t	t

\wedge	f	t
f	f	f
t	f	t

b	\bar{b}
f	t
t	f

WHAT IS A BOOLEAN ALGEBRA?

- $(B, +, *, \bar{}, \mathbf{0}, \mathbf{1})$
- Binary Operator $+$ (Sum)
- Binary Operator $*$ (Product)
- Unary Operator $\bar{}$ (Complement)
- $\mathbf{0}$ (Identity): $\forall b \in B, b + \mathbf{0} = \mathbf{0} + b = b$
- $\mathbf{1}$ (Identity): $\forall b \in B, b * \mathbf{1} = \mathbf{1} * b = b$
- $+$ and $*$ are Associative
- $+$ and $*$ are Commutative
- $+$ Distributes over $*$
- $*$ Distributes over $+$
- $\forall b \in B, b + \bar{b} = \mathbf{1}$
- $\forall b \in B, b * \bar{b} = \mathbf{0}$

BOOLEAN ALGEBRA PROPERTIES

- Dual - Change Binary operators and Identities
- The dual of a Boolean algebra theorem is a Boolean algebra theorem
- The identity elements are unique
- The complement of an element is unique
- $\forall b \in B, b + b = b$ and $b * b = b$
- $\forall b \in B, \mathbf{1} + b = \mathbf{1}$ and $\mathbf{0} * b = \mathbf{0}$
- Absorption: $\forall b_1, b_2 \in B, b_1 + (b_1 * b_2) = b_1$
and $b_1 * (b_1 + b_2) = b_1$
- Involution: $\forall b \in B, \overline{\overline{b}} = b$
- DeMorgan's Laws: $\forall b_1, b_2 \in B, \overline{(b_1 + b_2)} = \overline{b_1} * \overline{b_2}$
and $\overline{(b_1 * b_2)} = \overline{b_1} + \overline{b_2}$

BOOLEAN EXPRESSIONS

- Given $(B, +, *, \bar{}, \mathbf{0}, \mathbf{1})$, a Boolean Variable is a variable over the set B
- A Literal is a Boolean variable, x , or its complement \bar{x}
- Boolean Expressions
 - Identity Elements **0**, **1**
 - Boolean Variables x_1, x_2, \dots, x_n
 - $(X + Y), (X * Y), \overline{X}$ where X and Y are Boolean Expressions
- Two Boolean expressions are equivalent (equal) if one can be obtained from the other by a finite sequence of applications of the Boolean algebra axioms

GENERAL BOOLEAN ALGEBRAS

- The Power Set for any finite set can be used to define a Boolean Algebra
- Example: Look at $A=\{x,y\}$
- The cardinality of a finite Boolean algebra is a power of 2
- In Computer Science we are primarily interested in the Boolean algebra of cardinality 2

Boolean Algebra on Two Variables

Boolean Algebra on Four Variables

BOOLEAN FUNCTIONS

- A Boolean Function $f: B^n \rightarrow B$ such that $f(x_1, x_2, \dots, x_n)$ is a Boolean expression
- Examples: $f_1(x_1, x_2) = x_1 + x_2$; $f_2(x_1, x_2, x_3) = \overline{x_1} * (x_2 + \overline{x_3})$
- Notation: Use x_1x_2 to mean $x_1 * x_2$
- Precedence Order: Complement, Product, Sum, Left to Right

BOOLEAN FUNCTION ON ONE VARIABLE

- $f(x)=\mathbf{1}$, $g(x)=x$
- Examine $f(x)+f(x)$, $f(x)+g(x)$, $f(x)g(x)$, $g(x)+g(x)$,
 $g(x)g(x)$, $\overline{f(x)}$, $\overline{g(x)}$, $\overline{f(x)}$
- Any function $f(x)$ can be written as $f(x)=f(\mathbf{0})x+f(\mathbf{1})\bar{x}$.
- $\mathbf{0}x+\mathbf{0}\bar{x} = \mathbf{0}$
- $\mathbf{0}x+\mathbf{1}\bar{x} = \bar{x}$
- $\mathbf{1}x+\mathbf{0}\bar{x} = x$
- $\mathbf{1}x+\mathbf{1}\bar{x} = x+\bar{x} = \mathbf{1}$

MINTERMS & MAXTERMS

- One variable: x, \bar{x}
- Two variables: $x_1x_2, x_1\bar{x}_2, \bar{x}_1x_2, \bar{x}_1\bar{x}_2$
- Minterm on n variables x_1, x_2, \dots, x_n is a Boolean expression which has the form of the product of each Boolean variable or its complement.
- Two minterms on one variable, four on two, eight on 3, and 2^n on n .
- Notation: $x^0 = \bar{x}$ and $x^1 = x$.
- $m_{e_1e_2\dots e_n} = x_1^{e_1}x_2^{e_2}\dots x_n^{e_n}$
- Examples: $m_{11} = x_1x_2, m_{10} = x_1\bar{x}_2, m_{01} = \bar{x}_1x_2, m_{00} = \bar{x}_1\bar{x}_2$
- Maxterm on n variables x_1, x_2, \dots, x_n is a Boolean expression which has the form of the sum of each Boolean variable or its complement.
- Examples: $M_{11} = x_1 + x_2, M_{10} = x_1 + \bar{x}_2, m_{01} = \bar{x}_1 + x_2, m_{00} = \overline{x_1 + x_2}$

CANONICAL FORMS OF BOOLEAN EXPRESSIONS

- Any Boolean expression can be written as the sum of minterms (or product of maxterms)
- Disjunctive Normal Form - Boolean expression written as sum of minterms
- Conjunctive Normal Form - Boolean expression written as product of maxterms
- We will normally use DNF
- To convert a Boolean expression, $f(x_1, x_2, \dots, x_n)$ into DNF, we need to determine the values for the constant prefixes (That is the $f(e_1, e_2, \dots, e_n)$).

Convert $x_1 + x_2$ into DNF

CONVERTING TO CANONICAL FORM

- Algorithm:
 1. Create table showing values of $f(e_1, \dots, e_n)$
 2. Rewrite f using the output values shown as the prefixes for the corresponding terms

- Example:

e_1	e_2	$f(e_1, e_2)$
0	0	0
0	1	0
1	0	1
1	1	1

BOOLEAN ALGEBRA IN COMPUTER SCIENCE

- Switching Circuits
- Logic Networks
- Karnaugh Maps

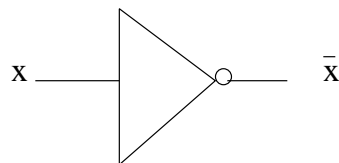
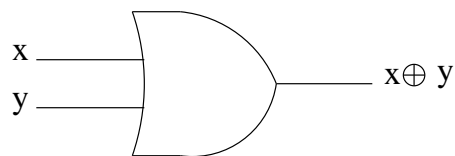
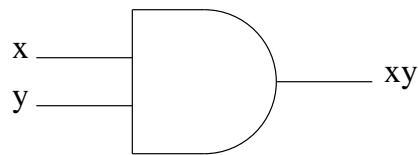
Design an adder

SWITCHING CIRCUITS

- Two-state Device: On or Off (1 or 0; T or F)
- Switch - Open (No current) or Closed (Current)
- Switches Can be combined:
 - Parallel - Boolean algebra +
 - Series - Boolean algebra *
 - If S denotes a switch then \bar{S} denotes a switch which always has the opposite state
- Switching functions can be created to represent any Boolean function
- Switches can be mechanical or electronic

LOGIC NETWORKS

- Switching circuits either conduct (switch closed) or do not conduct (switch open) electricity. For example voltage-operated circuits could define **1** as a signal with ≥ 3 volts and **0** with 0 volts and an acceptable tolerance
- Gate - hardware that implements logic operations of AND, OR, NOT (inverter)
- Linking gates together creates a logic network

**Logic Network Worksheet**