Software Quality Engineering:

Testing, Quality Assurance, and Quantifiable Improvement

Jeff Tian, tian@engr.smu.edu www.engr.smu.edu/~tian/SQEbook

Chapter 15. Formal Verification

- General idea and approaches
- Axiomatic verification
- Other approaches
- Summary and Perspectives

QA Alternatives

- Defect and QA:
 - Defect: error/fault/failure.
 - ▷ Defect prevention/removal/containment.
 - Map to major QA activities
- Defect prevention:
 Error source removal & error blocking
- Defect removal: Inspection/testing/etc.
- Defect containment: Fault tolerance and failure containment (safety assurance)
- Special case (this chapter): formal verification (& formal specification)

QA and Formal Verification

- Formal methods = formal specification
 - + formal verification
- Formal specification (FS):
 - > As part of defect prevention

 - ▷ Briefly covered as related to FV.
- Formal verification (FV):
 - As part of QA, but focus on positive: "Prove absence of fault"
 - ▶ People intensive
 - Several commonly used approaches
 - Chapter 15 focus on basic ideas

Formal Specification: Ideas

- Formal specification:
 - Correctness focus
 - Different levels of details
 - > 3Cs: complete, clear, consistent
- Descriptive formal specifications:

 - Math functions
 - Notations and language support:Z, VDM, etc.
- Behavioral formal specifications:
 FSM, Petri-Net, etc.

Formal Verification: Ideas

- "Testing shows the presence of errors, not their absence." — Dijkstra
- Formal verification: proof of correctness

 - > Axioms for components or functional units

 - > Development and verification together
- Other related approaches:
 - ▷ Semi-formal verification
 - ▶ Model checking
 - ▷ Inspection for correctness

Formal Verification Basics

- Basic approaches:
 - ⊳ Floyd/Hoare axiomatic
 - ▷ Dijkstra/Gries weakest precond. (WP)
 - Mills' prog calculus/functional approach
- Basis for verification:

 - ▶ mathematical function (Mills)
 - b other formalisms
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- Procedures/steps used:
 - bottom-up (axiomatic)
 - backward chaining (WP)
 - ⊳ forward composition (Mills), etc.

Object and General Approach

- Basic block: statements
 - block (begin/end)
- Formal verification
 - > rules for above units

Axiomatic Approach

- Floyd axioms/flowchart
 - > Annotation on flowchart
 - ▶ Logical relations
 - Verification using logic
- Hoare axioms/formalization
 - ▷ Pre/Post conditions

 - ▷ Invariants (loops, functions)

 - ⊳ Focus of Chapter 15

Axiomatic Correctness

Notations

- \triangleright Statements: S_i
- ▶ Logical conditions: {P} etc.
- Schema: {P} S {Q}

conditions or schemas conclusion

Axioms:

- Schema for assignment

- Loop invariant

Axiomatic Approach: Formal Specs

- Formal specification:

 - ▷ Pre-/post-conditions.
 - Pair as specifications at different levels of granularity.
- Example specification for a segment:
 - \triangleright Input/output variables: x, y.
 - \triangleright Pre-/post-conditions: P, Q.
 - \triangleright Pre-condition: non-negative input $\{P \equiv x \geq 0\}$
 - \triangleright Post-condition: square root computed $\{Q \equiv y = \sqrt{x}\}.$

Axiomatic Approach: Inference Rules

- Inference rules: Consequence axioms
 - ▶ Logical implications and deductions.
 - ⊳ Flexibility for different pre-/post-cond.
- Consequence 1: relaxing post-condition

Axiom A1 :
$$\frac{\{P\}S\{R\},\ \{R\}\Rightarrow\{Q\}}{\{P\}S\{Q\}}$$

Consequence 2: more strict pre-condition

Axiom A2:
$$\frac{\{P\}\Rightarrow\{R\},\ \{R\}S\{Q\}}{\{P\}S\{Q\}}$$

Compare to WP (later).

Axiomatic Approach: Axioms

- Assignment schema:
 - \triangleright Axiom A3: $\{P_x^y\}y \leftarrow x\{P\}$
 - \triangleright where $\{P_x^y\}$ is derived from P with all free occurrence of y replaced by x.
 - \triangleright Example: $b \leftarrow b w$ with
 - post-condition $b \ge 0$ (maintaining non-negative balance)
 - pre-condition is then $b-w\geq 0$ or $b\geq w$, sufficient fund for withdraw.
- Axiom A4. Sequential concatenation:

$$\frac{\{P\}S_1\{Q\},\ \{Q\}S_2\{R\}}{\{P\}S_1;S_2\{R\}}$$

Used to build bottom-up proofs.

Axiomatic Approach: Axioms

- Conditional axioms.
- Conditional 1, if-then-else (Axiom A5):

$$\frac{\{P \land B\}S_1\{Q\}, \{P \land \neg B\}S_2\{Q\}}{\{P\} \text{ if } B \text{ then } S_1 \text{ else } S_2 \{Q\}}$$

• Conditional 2, empty else (Axiom A6):

$$\frac{\{P \land B\}S\{Q\}, \ \{P \land \neg B\} \Rightarrow \{Q\}}{\{P\} \text{ if } B \text{ then } S \ \{Q\}}$$

Axiomatic Approach: Axioms

- Loop type: while cond do something
- Loop axiom (Axiom A7):

$$\frac{\{P \wedge B\}S\{P\}}{\{P\} \text{ while } B \text{ do } S \text{ } \{P \wedge \neg B\}}$$

- Specialized techniques for loops:
 - \triangleright Loop invariant: P (often labeled I)
 - ▶ How to select loop invariant?
 - ▶ Proof of basic loop: Axiom A7.
- Loop termination verification:
 - $\triangleright P$ positive within a loop
 - $\triangleright P_i > P_{i+1}$

Axiomatic Proofs

- Given: program, pre/post-conditions
- Basic proof procedure:
 - > Add annotations in between statements.
 - Apply axioms to individual statements using assignment schema (A3).
 - ⊳ Simple composition (concatenation, A4).
 - ▶ More complex composition:
 - if-then-else (A5) and if-then (A6)
 - loop axiom (A7): often the focus.
 - Consequence rules (A1 and A2) as connectors mixed with the above.
- General proof focuses:
 - Loop termination and invariants

 - Use hierarchical (stepwise abstraction)
 structure as guide for different parts
 (top-down guide bottom-up procedure)

Sample Axiomatic Proof

• Factorial function: Fig 15.1

• Sample axiomatic proof (pp.257-259)

Sample Axiomatic Proof

- Key to the proof: loop;
 other steps fairly straightforward.
- Loop invariant I development:
 - $\triangleright y$ holds partial results $\left(y = \frac{n!}{i!}\right)$.
 - \triangleright connection with loop condition i > 1.

$$\Rightarrow I \equiv \left(y = \frac{n!}{i!}\right) \land (i \ge 1).$$

 \triangleright resulting in post-condition after loop $I \land \neg B \equiv (y = n!)$

Observation from sample proof:
 proof much longer than the program itself

Axiomatic Proofs

- General observations:
 - Many steps involved
 - ▶ Length of proof: An order of magnitude longer than the program
 - Difficulty with loops
- Larger/more complex programs:
 - Many elements and (nested!) loops
 - ⇒ interaction, coordination
 - Arrays and functions/procedures
 - ⇒ more complicated schemas/axioms
 - Much harder.
 - Selective verification ideas?See Chapter 16, safety assurance part.

WP Approach

- Dijkstra/Gries approach:
 - \triangleright Weakest preconditions: wp(S,Q).
 - Dijkstra model: Predicate transforms.
 - ▷ Gries "Science of Programming" book.
- Similarity to axiomatic approach:
 - ▶ Logic based, same annotations.
 - ▷ Similar units (axioms).
 - $\triangleright \{P\}S\{Q\}$ interpreted as $P \Rightarrow wp(S,Q)$.
- Different procedures:

 - Backward chaining of WPs

Functional Approach

- Functional approach
 - ▶ Mills' program calculus
 - > Symbolic execution extensively used.
- Functional approach elements
 - Mills box notation
 - Basic function associated with individual statements
 - > Compositional rules

 - ▷ Comparison with Dijkstra's wp

Functional Approach: Symbolic Execution

Symbolic execution (Table 15.1, p.261) for:

if
$$x \geq 0$$
 then $y \leftarrow x$ else $y \leftarrow -x$

Trace 1:

| part | condition | X | У |
|------------------|-----------|---|---|
| if $x > 0$ | x > 0 | | |
| $y \leftarrow x$ | | | X |

Trace 2:

| part | condition | X | У |
|-------------------|------------|---|----|
| if $x > 0$ | $x \leq 0$ | | |
| $y \leftarrow -x$ | | | -x |

- Both traces used in verification
 - details in Mills et al. (1987a)

Formal Verification: Limitations

- Seven myths (Zelkowitz, 1993):
 - > FM guarantee that software is perfect.
 - ▶ They work by proving correctness.
 - Only highly critical system benefits.
 - > FM involve complex mathematics.
 - → FM increase cost of development.
 - ▶ They are incomprehensible to client.
 - ⊳ Nobody uses them for real projects.
- Refutation/discussion (Zelkowitz, 1993)
- However, some quantified validity
 - \Rightarrow alternative FV methods.

Other Models/Approaches

- Making FV more easily/widely usable.
- Other models for formal verification:
 - > State machines and model checking.
 - ▶ Algebraic data spec/verification.
 - ▶ Petri nets, etc.
- General assessment

 - ▶ More advantages & reduced limitations.
 - ▶ Formal analysis vs. verification.
 - ▶ May lead to additional automation.

 - > Adaptation and semi-formal methods.

Formal Verification: Other

- Algebraic specification/verification:
 - > Specify and verify data properties
 - ▶ Behavior specification
 - ▶ Base case

 - Domain/behavior mapping
- Stack example
 - ▷ newstack
 - ▷ push
 - ⊳ pop
 - Canonical form

Formal Verification: Other

Model checking:

- Behavioral specification via FSMs.
- > *Proposition*: property of interest expressed as a suitable formula.
- Model checker: algorithm/program to check proposition validity.
 - Proof: positive result.
 - Counterexample: negative result.

Other approaches and discussions:

- ▶ Algorithm analysis.
- ▶ Petri-net modeling and analysis.
- ▶ Tabular/semi-formal method.
- ▶ Formal inspection based.
- \triangleright Limited aspects \Rightarrow easier to perform.

FM: Applications

- What can be formally verified:
 - ▶ Program code.
 - ▶ Formal design, documentation, etc.
 - ▶ Protocols: timing properties
 - deadlock/starvation/etc.
 - ▶ Hardware verification.
 - Distributed program verification.
 - ▷ Connected to software process.
- Stepwise refinement/verification process:
 - ▷ Design and verification together.
 - Different levels of abstraction.
 - ▷ e.g., UNITY system.

Application in Software Safety

- Leveson approach (Chapter 16)

 - Driven by hazard analysis
 - Distributed over development phases
 - ▶ Which FM? as appropriate.
- Other applications:
 - Cleanroom:combination with statistical testing

Formal Verification: Summary

• Basic features:

- Axioms/rules for all language features
- ▶ Ignore some practical issues:
 Size, capacity, side effects, etc.?
- ▷ Develop invariants: key, but hard.

General assessment:

- Difficult, even on small programs
- Very hard to scale up
- ▷ Inappropriate to non-math. problems
- ▶ Hard to automate
 - manual process ⇒ errors↑
- Worthwhile for critical applications
- Comparison to other QA: Chapter 17.